## LECTURE - 24 <br> Semaphores and Bounded Buffer

## Semaphores

- Semaphore is a type of generalized lock
- Defined by Dijkstra in the last 60s
- Main synchronization primitives used in UNIX
- Consist of a positive integer value
- Two operations
$\bullet \mathbf{P (})$ : an atomic operation that waits for semaphore to become positive, then decrement it by 1
- V(): an atomic operation that increments semaphore by 1 and wakes up a waiting thread at $P()$, if any.


## Semaphores vs. Integers

- No negative values
- Only operations are P() and V()
- Cannot read or write semaphore values
- Except at the initialization times
- Operations are atomic
- Two P() calls cannot decrement the value below zero
- A sleeping thread at P() cannot miss a wakeup from V()


## Binary Semaphores

- A binary semaphore is initialized to 1
- P() waits until the value is 1
- Then set it to 0
$\diamond$ V() sets the value to 1
- Wakes up a thread waiting at $P($ ), if any


## Two Uses of Semaphores

1. Mutual exclusion

- Lock was designed to do this
lock->acquire() ;
// critical section
lock->release() ;


## Two Uses of Semaphores

1. Mutual exclusion
2. The lock function can be realized with a binary semaphore: semaphore subsumes lock.

- Semaphore has an initial value of 1
- P() is called before a critical section
$\bullet$ V() is called after the critical section
semaphore litter_box $=1$;
P (litter boz) ;
/ / critical section
V(1itter_boz) ;


## Two Uses of Semaphores

1. Mutual exclusion

- Semaphore has an initial value of 1
- P() is called before a critical section
- V() is called after the critical section
semaphore litter_box $=1$;
P (litter_box) ;
litter_box = 1
// critical section
V(litter_box) ;


## Two Uses of Semaphores

1. Mutual exclusion

- Semaphore has an initial value of 1
- $P()$ is called before a critical section
- V() is called after the critical section
semaphore litter_box = 1;
P(litter_box) ; // purrx... litter_box=1 $\rightarrow 0$
// critical section
v(litter_box) ;


## Two Uses of Semaphores

1. Mutual exclusion

- Semaphore has an initial value of 1
- P() is called before a critical section
-V() is called after the critical section
semaphore litter_box = 1;
P (litter_box) ;

// critical section
V(litter_box) ;


## Two Uses of Semaphores

1. Mutual exclusion

- Semaphore has an initial value of 1
- $P($ ) is called before a critical section
- V() is called after the critical section
semaphore litter_box = 1;
P(litter_box) ; // meow... \& litter_box = 0
// critical section
v(litter_box) ;


## Two Uses of Semaphores

1. Mutual exclusion

- Semaphore has an initial value of 1
- $P($ ) is called before a critical section
- V() is called after the critical section
semaphore litter_box = 1;
P(litter_box) ;

$$
\text { litter_box = } 0 \rightarrow 1
$$

// critical section
v(litter_box) ;

## Two Uses of Semaphores

2. Synchronization: Enforcing some order between threads


## Two Uses of Semaphores

## 2. Synchronization

- Semaphore usually has an initial value of 0

```
semaphore wait_left = 0;
semaphore wait_right = 0;
Ieft_Paw() {
    slide_left();
    V(wait_left);
    P(wait_right);
    slide_right();
}
```

```
Right_Paw() {
```

Right_Paw() {
P(wait_left);
P(wait_left);
slide_left();
slide_left();
slide_right();
slide_right();
V (wait_right) ;
V (wait_right) ;
}

```
}
```


## Two Uses of Semaphores

## 2. Scheduling

- Semaphore usually has an initial value of 0

```
semaphore wait_left = 0;
semaphore wait_right = 0;
```



```
Left_Paw() {
Right_Paw() {
    slide_left();
    V(wait_left);
    P(wait_right);
    slide_right();
}
}
```


## Two Uses of Semaphores

## 2. Scheduling

- Semaphore usually has an initial value of 0

```
semaphore wait_left = 0;
semaphore wait_right = 0;
```

```
Left_Paw() {
```

Left_Paw() {
slide_left();
slide_left();
V(wait_left);
V(wait_left);
P(wait_right);
P(wait_right);
slide_right();
slide_right();
}

```
```

Right_Paw() {

```
Right_Paw() {
    P(wait_left);
    P(wait_left);
    slide_left();
    slide_left();
    slide_right();
    slide_right();
    V(wait_right);
    V(wait_right);
}
```

}

```
wait_left = 0
wait_right = 0

\section*{Two Uses of Semaphores}

\section*{2. Scheduling}
- Semaphore usually has an initial value of 0
```

semaphore wait_left = 0;
semaphore wait_right = 0;

```
wait_left = 0
wait_right \(=0\)
```

Left_Paw() {
slide_left();
V(wait_left);
P(wait_right);
slide_right();
}

```
Right_Paw() \{
    P(wait_left) ;
    slide_left () ;
    slide_right () ;
    V(wait_right) ;

\section*{Two Uses of Semaphores}

\section*{2. Scheduling}
- Semaphore usually has an initial value of 0
```

semaphore wait_left = 0;
semaphore wait_right = 0;

```
wait_left = 0
wait_right \(=0\)
```

Ieft_Paw() {
slide_left();
V(wait_left);
P(wait_right);
slide_right();
}

```
```

Right_Paw() {
P(wait_left);
slide_left();
slide_right();
V (wait_right) ;
}

```

\section*{Two Uses of Semaphores}

\section*{2. Scheduling}
- Semaphore usually has an initial value of 0
```

semaphore wait_left = 0;
semaphore wait_right = 0;

```
wait_left \(=0 \rightarrow 1\)
wait_right \(=0\)
```

Ieft_Paw() {
slide_left();
V(wait_left);
P(wait_right);
slide_right();
}

```
```

Right_Paw() {
P(wait_left);
slide_left();
slide_right();
V (wait_right) ;
}

```

\section*{Two Uses of Semaphores}

\section*{2. Scheduling}
- Semaphore usually has an initial value of 0
```

semaphore wait_left = 0;
semaphore wait_right = 0;

```
wait_left \(=1 \rightarrow 0\)
wait_right \(=0\)
```

Ieft_Paw() {
Right_Paw() {
slide_left();
V(wait_left);
P(wait_right);
slide_right();
}
}

```

\section*{Two Uses of Semaphores}

\section*{2. Scheduling}
- Semaphore usually has an initial value of 0
```

semaphore wait_left = 0;
semaphore wait_right = 0;
slide_left();
V (wait_left);
P(wait_right);
slide_right();
}

Left_Paw () \{

```
Right_Paw() {
```

Right_Paw() {
P(wait_left);
P(wait_left);
slide_left();
slide_left();
slide_right();
slide_right();
V(wait_right);
V(wait_right);
}

```
```

}

```
```

    wait_left = 0
    wait_right = 0
    
## Two Uses of Semaphores

## 2. Scheduling

- Semaphore usually has an initial value of 0

```
semaphore wait_left = 0;
semaphore wait_right = 0;
```

wait_left = 0
wait_right $=0$

Left_Paw () \{
slide_left();
V (wait_left) ;
P(wait_right) ;
slide_right () ;


```
Right_Paw() {
    P(wait_left);
    slide_left();
    slide_right();
    V (wait_right) ;
}
```


## Two Uses of Semaphores

## 2. Scheduling

- Semaphore usually has an initial value of 0

```
semaphore wait_left = 0;
semaphore wait_right = 0;
```

wait_left = 0
wait_right $=0$

```
Ieft_Paw() {
    slide_left();
    V(wait_left);
    P(wait_right);
    slide_right();
}
```

```
Right_Paw() {
    P(wait_left);
    slide_left();
    slide_right();
    V (wait_right) ;
}
```


## Two Uses of Semaphores

## 2. Scheduling

- Semaphore usually has an initial value of 0

```
semaphore wait_left = 0;
semaphore wait_right = 0;
```

wait_left = 0
wait_right $=0 \rightarrow 1$

Left_Paw () \{
slide_left();
V(wait_left);
P(wait_right) ;
slide_right () ; \}
\}

## Two Uses of Semaphores

## 2. Scheduling

- Semaphore usually has an initial value of 0

```
semaphore wait_left = 0;
semaphore wait_right = 0;
```

wait_left = 0
wait_right $=1 \rightarrow 0$

```
Ieft_Paw() {
    slide_left();
    V (wait_left);
    P(wait_right) ;
    slide_right();
}
    Right_Paw() {
    P(wait_left);
    slide_left();
    slide_right();
    V (wait_right) ;
        }
```


## Two Uses of Semaphores

## 2. Scheduling

- Semaphore usually has an initial value of 0

```
semaphore wait_left = 0;
semaphore wait_right = 0;
```

wait_left = 0
wait_right $=0$

Left_Paw () \{
slide_left () ;
V(wait_left);
P(wait_right) ;
slide_right () ;
\}

```
Right_Paw() {
    P(wait_left);
    slide_left();
    slide_right();
    V (wait_right) ;
}
```


## Two Uses of Semaphores

## 2. Synchronization

- Semaphore usually has an initial value of 0

```
semaphore s1 = 0;
semaphore s2 = 0;
A() {
```

```
B() {
```

B() {
P(s1);
P(s1);
read (x) ;
read (x) ;
write(y);
write(y);
V(s2);
V(s2);
}

```
}
```

